Local equivalence of stabilizer states

a graphical characterisation

Nathan Claudet and Simon Perdrix

QIP 2025 - 27/02/25 arXiv:2409.20183





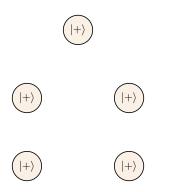






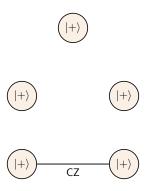


Graph states, local unitary equivalence, local Clifford equivalence & local complementation



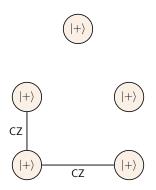
¹Edges do not have a direction.

²No multiples edges and no loops.



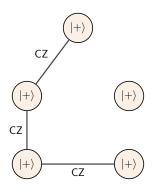
¹Edges do not have a direction.

²No multiples edges and no loops.



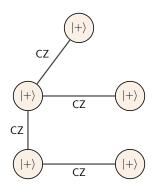
¹Edges do not have a direction.

²No multiples edges and no loops.



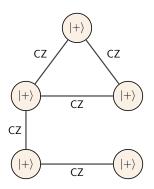
¹Edges do not have a direction.

²No multiples edges and no loops.



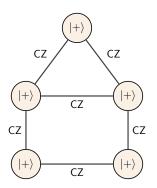
¹Edges do not have a direction.

²No multiples edges and no loops.



¹Edges do not have a direction.

²No multiples edges and no loops.

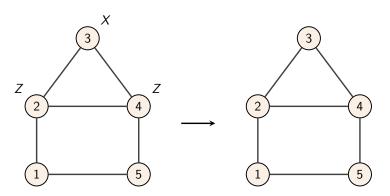


¹Edges do not have a direction.

²No multiples edges and no loops.

Stabilizer states

Graph states are a subfamily of stabilizer states because for each vertex u, applying X on u and applying Z on the neighbours of u leaves the graph state invariant.



4

Graph states are useful entangled resources (MBQC, error correction...).

ightarrow It is a fundamental problem to know whether two graph states have the same entanglement, i.e. the graph states are related by SLOCC.

Graph states are useful entangled resources (MBQC, error correction...).

 \rightarrow It is a fundamental problem to know whether two graph states have the same entanglement, i.e. the graph states are related by SLOCC.

Theorem (Van den Nest, Dehaene, De Moor, 2004)

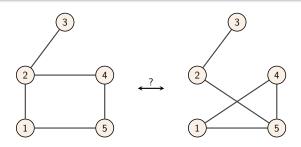
Two graph states are SLOCC-equivalent iff they are local unitary equivalent (or LU-equivalent), i.e. they are related by a tensor product of single-qubit unitaries.

Graph states are useful entangled resources (MBQC, error correction...).

 \to It is a fundamental problem to know whether two graph states have the same entanglement, i.e. the graph states are related by SLOCC.

Theorem (Van den Nest, Dehaene, De Moor, 2004)

Two graph states are SLOCC-equivalent iff they are local unitary equivalent (or LU-equivalent), i.e. they are related by a tensor product of single-qubit unitaries.

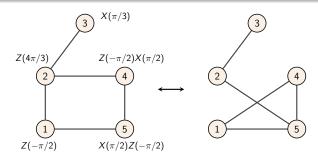


Graph states are useful entangled resources (MBQC, error correction...).

 \to It is a fundamental problem to know whether two graph states have the same entanglement, i.e. the graph states are related by SLOCC.

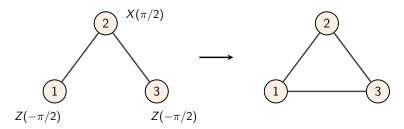
Theorem (Van den Nest, Dehaene, De Moor, 2004)

Two graph states are SLOCC-equivalent iff they are local unitary equivalent (or LU-equivalent), i.e. they are related by a tensor product of single-qubit unitaries.



An easier subproblem: local Clifford equivalence

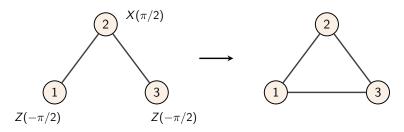
Two graph states are said **local Clifford equivalent** (or LC-equivalent) if they are related by unitaries in the local Clifford group.



6

An easier subproblem: local Clifford equivalence

Two graph states are said **local Clifford equivalent** (or LC-equivalent) if they are related by unitaries in the local Clifford group.



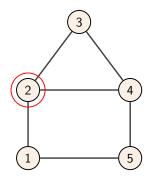
Theorem (Van den Nest, Dehaene, De Moor, 2004)

Two graph states are local Clifford equivalent iff the two corresponding graphs are related by **local complementations**.

Local complementation

Definition (Kotzig, 1966)

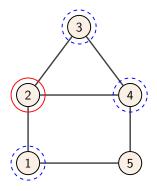
A local complementation on a vertex u consists in complementing the (open) neighbourhood of u.



Local complementation

Definition (Kotzig, 1966)

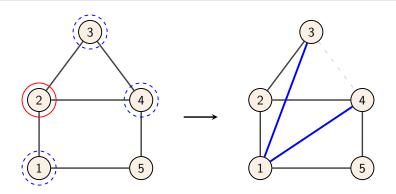
A local complementation on a vertex u consists in complementing the (open) neighbourhood of u.



Local complementation

Definition (Kotzig, 1966)

A local complementation on a vertex u consists in complementing the (open) neighbourhood of u.



7

Algorithmic aspect of local Clifford equivalence

There exists an efficient algorithm (Bouchet, 1991) to recognise whether two graphs are related by local complementations, implying an efficient algorithm to recognise whether two graph states are local Clifford equivalent.

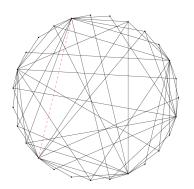
$LU \neq LC$

Unfortunately, LU \neq LC, i.e. local Clifford equivalence and local unitary equivalence do not coincide.

Ç

$LU \neq LC$

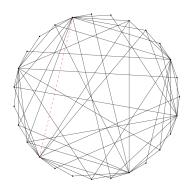
Unfortunately, LU \neq LC, i.e. local Clifford equivalence and local unitary equivalence do not coincide. \leftarrow 27-qubit pair of graph states that are local unitary equivalent but not local Clifford equivalent (Ji et al. 2008).



9

$LU \neq LC$

Unfortunately, LU \neq LC, i.e. local Clifford equivalence and local unitary equivalence do not coincide. \leftarrow 27-qubit pair of graph states that are local unitary equivalent but not local Clifford equivalent (Ji et al. 2008).



Consequence: local complementation does **not** capture the local unitary equivalence of graph states.

9

They are some known families of graph states for which LU=LC i.e. local complementation captures local unitary equivalence:

• Graph states over at most 8 qubits (Cabello et al. 2009)

- Graph states over at most 8 qubits (Cabello et al. 2009)
- Complete graphs i.e. GHZ states (Van den Nest, Dehaene, De Moor, 2005)

- Graph states over at most 8 qubits (Cabello et al. 2009)
- Complete graphs i.e. GHZ states (Van den Nest, Dehaene, De Moor, 2005)
- Complete bipartite graphs (Tzitrin, 2018)

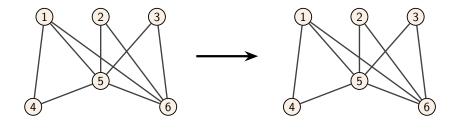
- Graph states over at most 8 qubits (Cabello et al. 2009)
- Complete graphs i.e. GHZ states (Van den Nest, Dehaene, De Moor, 2005)
- Complete bipartite graphs (Tzitrin, 2018)
- Graphs with no cycle of length 3 or 4 (Zeng et al. 2007)

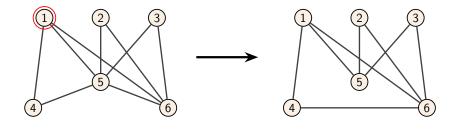
They are some known families of graph states for which LU=LC i.e. local complementation captures local unitary equivalence:

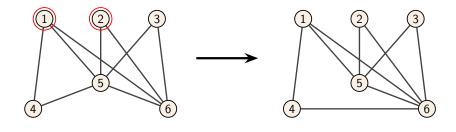
- Graph states over at most 8 qubits (Cabello et al. 2009)
- Complete graphs i.e. GHZ states (Van den Nest, Dehaene, De Moor, 2005)
- Complete bipartite graphs (Tzitrin, 2018)
- Graphs with no cycle of length 3 or 4 (Zeng et al. 2007)

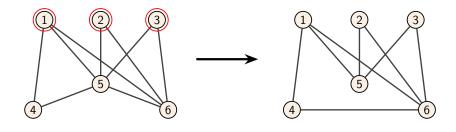
But what about local unitary equivalence for **any** graph? Can we construct a graphical characterisation?

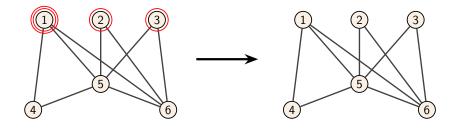
Generalising local complementation to capture local unitary equivalence

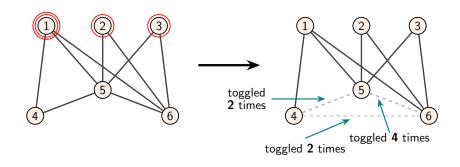






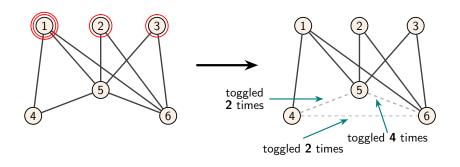






A refinement of idempotent local complementations

A sequence of local complementations may leave the graph invariant.

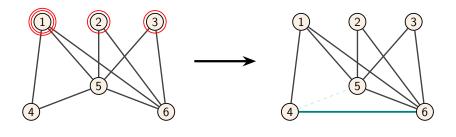


toggled 2 mod 4 times by the idempotent local complementations. (There are also some additional conditions on the edges for the 2-local complementation the be valid.)

A **2-local complementation** consists in toggling every edge that was

A refinement of idempotent local complementations

A sequence of local complementations may leave the graph invariant.



A **2-local complementation** consists in toggling every edge that was toggled 2 mod 4 times by the idempotent local complementations. (There are also some additional conditions on the edges for the 2-local complementation the be valid.)

r-local complementation

- 3-local complementation is a refinement of idempotent 2-local complementation, and so on...
- \rightarrow Infinite family of graphical operations parametrised by an integer r:

r-local complementations

1-local complementation = local complementation.

Main results

Theorem (this work)

Two graphs are related by r-local complementations iff the two corresponding graph states are related by local unitaries in the level r+1 of the Clifford hierarchy.

For r = 1, we recover local Clifford \Leftrightarrow local complementation.

Main results

Theorem (this work)

Two graphs are related by r-local complementations iff the two corresponding graph states are related by local unitaries in the level r+1 of the Clifford hierarchy.

For r = 1, we recover local Clifford \Leftrightarrow local complementation.

Theorem (this work)

Two graph states are local unitary equivalent iff the two corresponding graphs are related by r-local complementations for some r.

Main results

Theorem (this work)

Two graphs are related by r-local complementations iff the two corresponding graph states are related by local unitaries in the level r+1 of the Clifford hierarchy.

For r = 1, we recover local Clifford \Leftrightarrow local complementation.

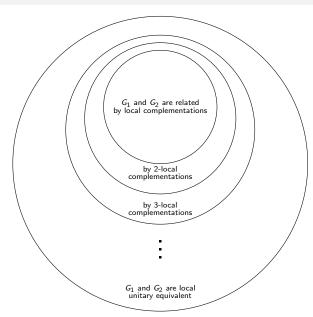
Theorem (this work)

Two graph states are local unitary equivalent iff the two corresponding graphs are related by r-local complementations for some r.

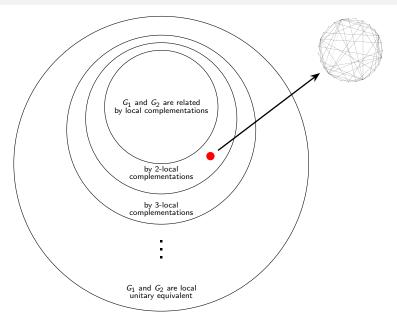
Corollary

If two graph states are local unitary equivalent, the local unitaries can be chosen to be in the Clifford hierarchy.

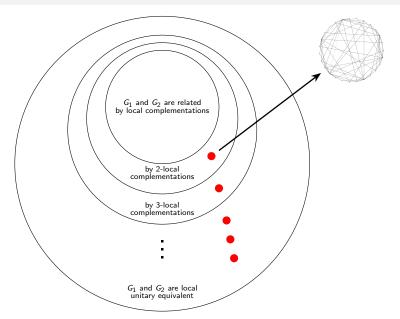
An infinite hierarchy of local equivalences



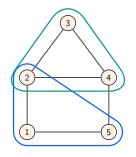
An infinite hierarchy of local equivalences



An infinite hierarchy of local equivalences



Proof sketch: Minimal local set

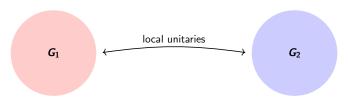


Minimal local sets are subsets of vertices that are invariant by local unitary equivalence and carry information on the possible local unitaries that maps graph states to other graph states.

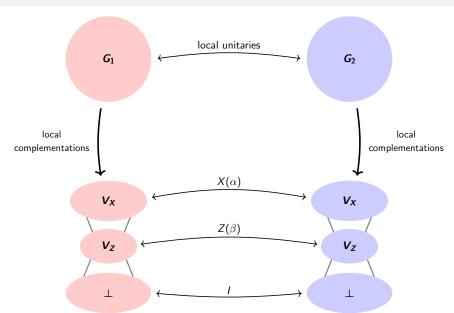
Theorem (C, Perdrix, 2024)

Each vertex of a graph is covered by at least one minimal local set.

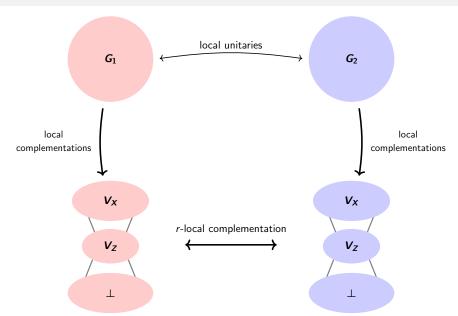
Proof sketch: Standard form for graph states



Proof sketch: Standard form for graph states

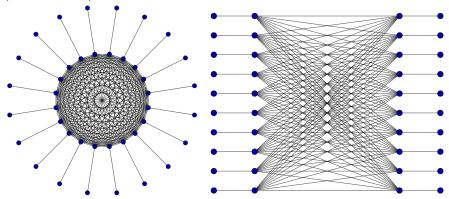


Proof sketch: Standard form for graph states



Application 1: local equivalence of repeater graph states

It was conjectured that LU=LC holds for some repeater graph states (Tzitrin, 2018). We showed that this is indeed the case.



Application 2: LU=LC for graph states up to 19 qubits

Previously: LU=LC for graph states up to 8 qubits, and there exists a 27-qubit pair for which LU \neq LC.

Application 2: LU=LC for graph states up to 19 qubits

Previously: LU=LC for graph states up to 8 qubits, and there exists a 27-qubit pair for which LU \neq LC.

Theorem (C, Perdrix, 2025)

LU=LC for graph states up to 19 qubits.

Application 3: A quasi-polynomial algorithm to decide local unitary equivalence

Previously: exponential algorithm for deciding local unitary equivalence of graph states (Burchardt, de Jong, Vandré, 2024).

Application 3: A quasi-polynomial algorithm to decide local unitary equivalence

Previously: exponential algorithm for deciding local unitary equivalence of graph states (Burchardt, de Jong, Vandré, 2024).

Theorem (C, Perdrix, 2025)

There exists an algorithm that decides whether two graph states are local unitary equivalent with runtime $n^{\log_2(n)+O(1)}$.

The generalised local complementation is a graph rule that completely captures the local unitary equivalence of graph states.

The generalised local complementation is a graph rule that completely captures the local unitary equivalence of graph states.

Open questions:

 Does there exist a counter-example to LU=LC between 20 and 26 qubits?

The generalised local complementation is a graph rule that completely captures the local unitary equivalence of graph states.

Open questions:

- Does there exist a counter-example to LU=LC between 20 and 26 qubits?
- Does there exist a polynomial-time algorithm for local unitary equivalence?

Thanks



arXiv:2409.20183