# Recognizing graph states with the same entanglement by looking at the underlying graphs

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# Graph states, local unitary equivalence, local Clifford equivalence & local complementation



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Graph states are a subfamilly of stabilizer states because for each vertex u, applying X on u and applying Z on the neighbours of u leaves the graph state invariant.



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## An easier subproblem: local Clifford equivalence

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Two graph states are LC-equivalent iff the two corresponding graphs are related by **local complementations**.

#### Definition (Kotzig, 1966)



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# Algorithmic aspect of local Clifford equivalence

#### Proposition (Bouchet, 1991)

There exists an efficient algorithm to decide if two graphs are related by local complementations.

 $\rightarrow$  algorithm to decide if two graph states are LC-equivalent.

# Quick history of the LU=LC conjecture

# The LU=LC conjecture

Formulated in the early 2000's.

#### Conjecture

LU=LC i.e. if two graph states are LU-equivalent (local unitaries) then they are LC-equivalent (local Clifford).

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#### Proposition (Gross, Van den Nest, 2007)

If two graph states are LU-equivalent, the local unitaries can be chosen to be of the form  $C_1Z(\theta)C_2$  where  $C_1$  and  $C_2$  are Clifford.

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However,  $LU \neq LC$ , i.e. local unitary equivalence and local Clifford equivalence do not coincide. $\rightarrow$  27-qubit pair of graph states that are LU-equivalent but not LC-equivalent (Ji et al. 2008).



## Another look at the 27-vertex counterexample

The 27-vertex counterexample is LC-equivalent to a prettier pair of graphs (Tsimakuridze, Gühne, 2017).



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But what about LU-equivalence for **any** graph ? Can we construct a graphical characterisation ?

# Generalizing local complementation to capture local unitary equivalence













A sequence of local complementations may leave the graph invariant.



A **2-local complementation** consists in toggling every edge that was toggled 2 mod 4 times by the idempotent local complementations. (There are also some additional conditions on the edges for the 2-local complementation the be valid.)

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# Example of a 2-local complementation



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3-local complementation is a refinement of idempotent 2-local complementation, and so on...

 $\rightarrow$  Infinite family of graphical operations parametrised by an integer r:

#### *r*-local complementations

1-local complementation = local complementation.

Recall: LC-equivalent  $\Leftrightarrow$  related by local unitaries in  $\langle H, Z(\pi/2) \rangle$ . Define: **LC**<sub>r</sub>-equivalent  $\Leftrightarrow$  related by local unitaries in  $\langle H, Z(\pi/2^r) \rangle$ . Recall: LC-equivalent  $\Leftrightarrow$  related by local unitaries in  $\langle H, Z(\pi/2) \rangle$ . Define: **LC**<sub>r</sub>-equivalent  $\Leftrightarrow$  related by local unitaries in  $\langle H, Z(\pi/2^r) \rangle$ .

#### Theorem (C, Perdrix, 2025)

Two graph states are  $LC_r$ -equivalent iff the two corresponding graphs are related by r-local complementations.

For r = 1, we recover local Clifford  $\Leftrightarrow$  local complementation.

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#### Theorem (<u>C</u>, Perdrix, 2025)

Two graph states are LU-equivalent iff the two corresponding graphs are related by r-local complementations for some r.

#### An infinite hierarchy of local equivalences



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# Proof that r-local complementation captures LU-equivalence

#### Definition (Odd neighbourhood)

Given a set of vertices D, the **odd neighbourhood**  $Odd_G(D)$  of D is the set of vertices that are neighbours of an odd number of vertices in D.





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#### Definition

A **local set** is a non-empty vertex set of the form  $L = D \cup Odd_G(D)$ . A **minimal local set** is a local set that is minimal by inclusion (i.e it doesn't strictly contain another local set).





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Minimal local sets carry information on the possible local unitaries that maps graph states to other graph states.

#### Minimal local sets cover any graph

Theorem (C, Perdrix, 2024)

Each vertex of a graph is covered by at least one minimal local set.



#### Proof sketch: Standard form for graph states



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# Application 1: a toolbox to prove LU=LC for classes of graphs
#### Proposition (C, Perdrix, 2025)

Suppose  $n \leq 2^{r+3} - 1$ . If two graph states over n qubits are LU-equivalent then they are  $LC_r$ -equivalent i.e.  $LU=LC_r$ .

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#### Corollary

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#### Proposition (C, Perdrix, 2025)

LU=LC for graph states up to 19 qubits.

# A graphical characterization for LU=LC

#### Proposition (C, Perdrix, 2024)

LU=LC for  $|G\rangle$  if there exists  $|G'\rangle$  such that:

- $|G\rangle$  and  $|G'\rangle$  are LC-equivalent;
- G' is in standard form;
- Any r-local complementation over V<sub>X</sub> can be implemented by usual local complementations.



### Example: LU=LC for repeater graph states

It was conjectured that LU=LC holds for some repeater graph states (Tzitrin, 2018). We showed that this is indeed the case.



# Application 2: a quasi-polynomial algorithm for LU-equivalence

Application 3: A quasi-polynomial algorithm to decide local unitary equivalence

Theorem (C, Perdrix, 2025)

There exists an algorithm that decides if two graph states are LU-equivalent with runtime  $n^{\log_2(n)+O(1)}$ .

## Step 1: Standard form



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# Step 2: Reduction to LC-equivalence with constraints



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# Summary

• An infinite strict hierarchy of local equivalence;

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- LU=LC for graph states up to 19 qubits;
- A quasi-polynomial algorithm for LU-equivalence.

• Does there exist a counter-example to the LU=LC conjecture between 20 and 26 qubits ?

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• Does there exist a polynomial-time algorithm for LU-equivalence ?

# Thanks



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